

## **A STUDY ON EDGES SET OF GRAPHS**

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### **ABSTRACT**

Diagram based component determination methodologies that can diminish the dimensionality of informational indexes while keeping up with the mathematical construction. To assess and think about the chose highlights, measures like grouping exactness, standardized shared data, changed common data, and the Jaccard file can be used. With informational collections of high component measurement and low example measurement, the five offered highlight choice calculations work rapidly. Be that as it may, regardless of the unassuming component measurement, it takes altogether more for informational indexes with a high example measurement. Dividing the information, utilizing designs preparing units, or investigating other figuring techniques reasonable for huge scope information should all be considered for viability. Subsequent to lessening the measure of beginning components, prototypical relapse models can be created. The nature of the models would then be able to be analyzed as far as forecast blunders, which might even demonstrate if the chose components ought to be thought of. A potential subsequent examination subject is tweaking chart-based component choices by redoing a few capacities and measures dependent on the informational collection's provisions and the issue solver's objectives. Elective techniques for producing the weight framework, gathering information focuses, and contrasting chosen highlights in split information, for instance, are captivating.

***KEY WORDS: Diagram, Edges, Set, Graphs, Captivating, Informational Indexes.***

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## 1. INTRODUCTION:

The study of sets loads two ways the first one is path given by logicians. They have analyzed theory in great detail and have formulated axioms for the subject each of their axioms expresses a property of sets that mathematicians commonly accept. The other goes up, onto the high lands of mathematics itself where these concepts are indispensable in almost all of pure mathematics as it is today here in this research, we introduce the ideas of set theory and establish the basic terminology and notation. The new thing which we shall give in this chapter is the concept of 'set of graphs'. The set of graphs will consist only simple graphs.

### *Set :*

A set is a well-defined collection of objects each of which satisfies a certain property such that it enables us to decide as to whether the given objects belong to that collection or not. Commonly we shall use capital letter A, B, ..... to denote the sets and small letters a, b, c..... to denote the objects of elements belonging to these sets.

If an objects  $x$  belongs to a set  $A$  we denote it by the notation  $x \in A$ .

If  $x$  does not belong to  $A$  we express it by the notation  $x \notin A$ .

We say that  $A$  is a subset of  $B$  if every element of  $A$  is also an element of  $B$  and are we express by  $A \subseteq B$ .

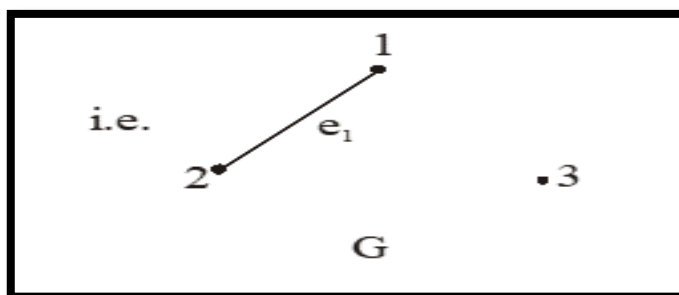
If  $A=B$  it is g true that both  $A \subseteq B$  and  $B \subseteq A$ . If  $A \subseteq B$  and  $B \subseteq A$ . If  $A \subseteq B$  and  $A$  is different from  $B$  we say that  $A$  is proper subset of  $B$  and we write  $A \subset B$ .

### 1.1 SET OF GRAPHS:

Let us assume a simple graph  $G (V, E)$ , where  $V$  is the set of vertices and  $E$  is the set of edges. In our course of study, we shall concern ourself only on the set of edges  $E$  and throughout the work it is called as the set of graphs. In a simple graph  $G (V, E)$ , if there are  $n$  edges then the set of graphs will contain  $2^n$  sub graphs with respect to the edge set i.e.,  $E$ .

Now let us take some examples to clear the concept of the set of graphs.

**Example:** Let us consider a simple graph  $G(V,E)$ , where  $V=\{1,2,3\}$  and the edge set  $E=\{e_1\}$ .



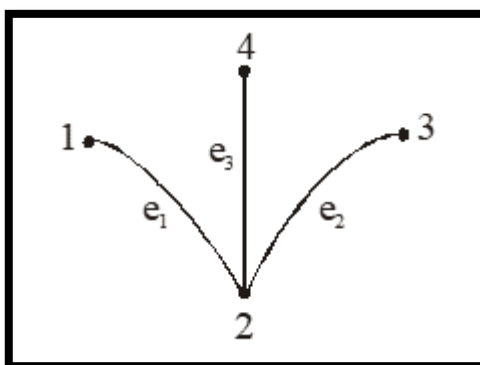
As there exists only one edge in the graph  $G$ , hence  $\exists 2^1$  elements in the set of graphs. thus the subset of  $E$  will be  $\phi$  &  $\{e_1\}$ .

Thus the elements in the set of graphs are  $\{\phi, \{e_1\}\}$ .



For the better understanding of the above given concept, let us take two more example.

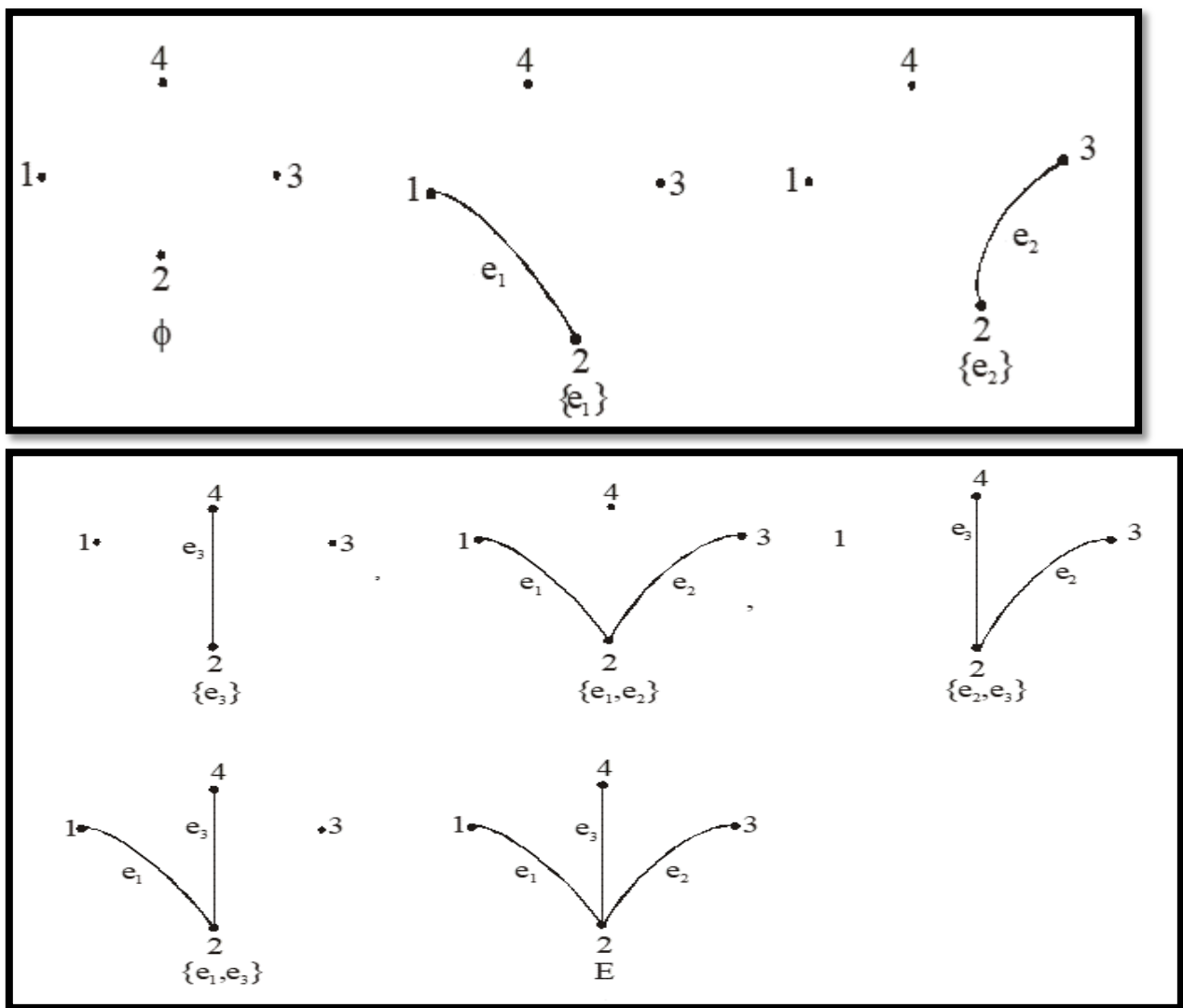
**Example :** Let us take a simple graph  $G(V,E)$ , where  $V = \{1,2,3,4\}$  and the edge set  $E = \{e_1, e_2, e_3\}$ .



The given graph G is defined over three edges. Thus the set of graphs will contain  $2^3$  i.e.. 8 sub graphs.

Set of graphs =  $\{\phi, E, \{e_1\}, \{e_2\}, \{e_3\}, \{e_1, e_2\}, \{e_1, e_3\}, \{e_2, e_3\}\}$ .

This set of graphs can be represented as follows:



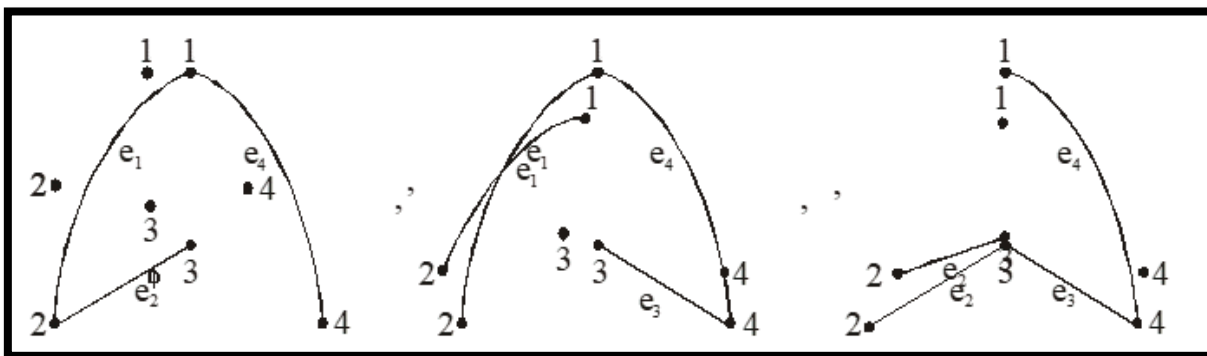
Here we can observe that null set and the edge set E always be member of the set of graphs.

**Example:** Similarly, we will find the set of graphs on the edge set of four edges.

Let us assume an edge set  $E = \{e_1, e_2, e_3, e_4\}$  on the simple graph  $G(V, E)$ .

As it is given that there exist four edges. Then by the above defined deaf. there will exist  $2^4$  sub graphs *i.e.*, 16 sub graphs.

Thus, define these elements of the set of graphs *i.e*



Now, we have come to a conclusion that the set of graphs can be defined on any simple graph like a any set of elements.

## 1.2 TYPES OF EDGE SET:

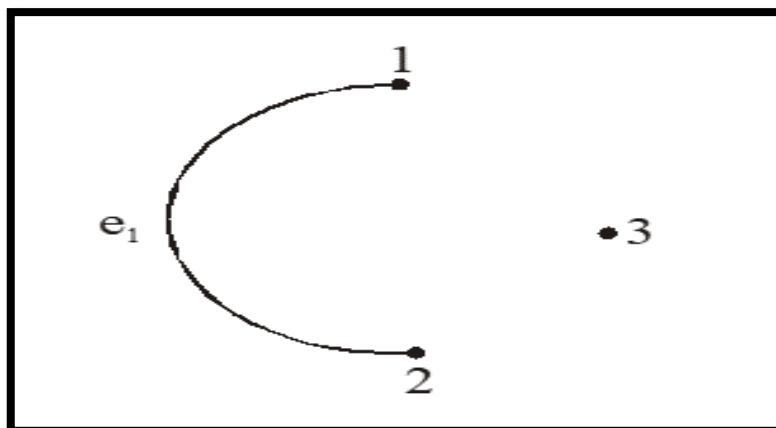
The set of edges have same types of sets as in case of simple sets of elements. Different types of edge set can express as follows:

### *Singleton Set:*

An edge set will be singleton set if it contains only one element *i.e.*, only on edge.

**Example :** let us consider a simple graph containing a single edge. Let the graph is  $g(V, E)$ , where  $n(V)=3$  simple edge. Let the graph is  $G(V, E)$ , where  $n(V) = 3$  and  $\eta(E) = 1$  *i.e.*  $V = \{1,2,3\}$  &  $E = \{e_1\}$ .

As the graph  $G$  has only one edge, so it is called singleton set irrespective of the fact that it contains their vertices.

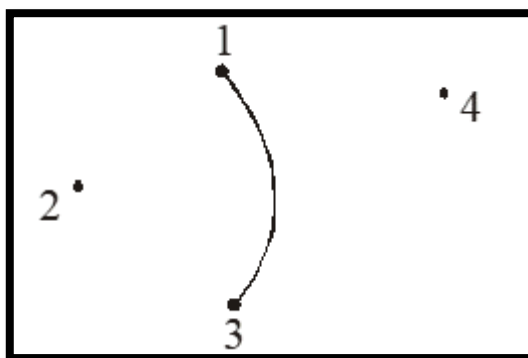


One more example of singleton set is as follows: -

**Example:-** Let the graph is  $G(V,E)$ , where no. of vertices are 4 and no. of edges is 1. i.e.

$V=\{1,2,3,4\}$ ,  $E=\{e_1\}$  because the graph  $G$  has only one edge, so it is called singleton set.

The graph is as follows:-



### **Null Set:**

Any edge set is said to be the null set if it contains no edge i.e. empty edge set of a simple graph  $G(V,E)$ . where  $\eta(E)=\phi$ .

**Example :** let us consider an example for better understanding.

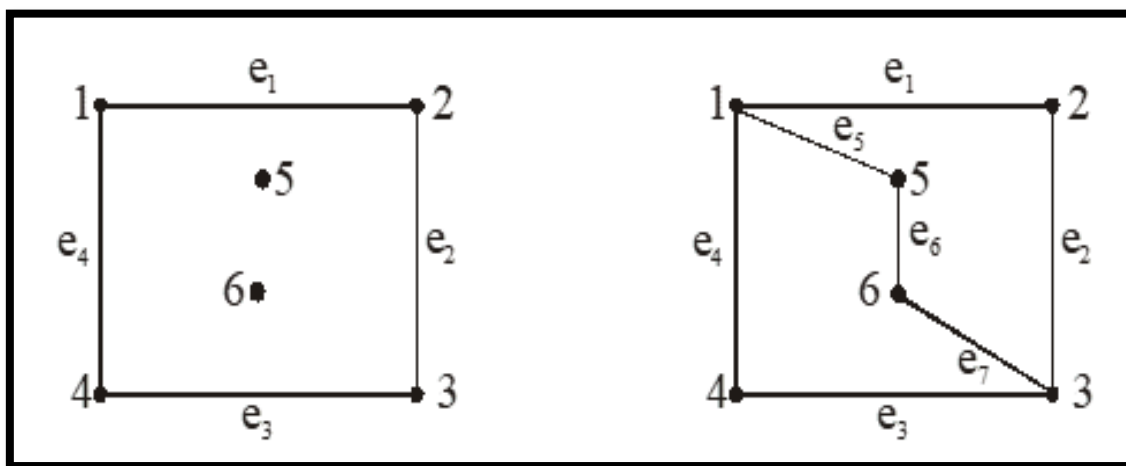
let the graph  $G(V,E)$  is a simple graph, where  $\eta(V)=3$  i.e.,  $V\{1,2,3\}$  and  $\eta(E)=\phi$  i.e.  $E\{\}$

**Subset or Superset:** Two edge set  $E_1$ , and  $E_2$ , are such that every edge of  $E_1$  belongs to the edge set  $E_2$ , then  $E_1$  is called the subset of  $E_2$ , &  $E_2$  is called the superset of  $E_1$ ,

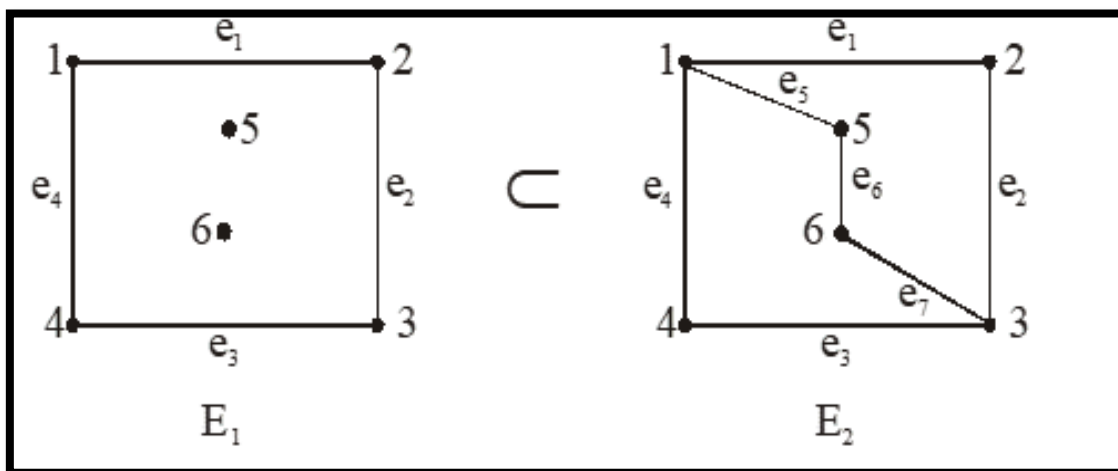
**Example:** Let us assume two edge set  $E_1$  and  $E_2$ , such that  $t_1 = \{e_1, e_2, e_3, e_4\}$  and the edge set  $E_2 = \{e_1, e_2, e_3, e_4, e_5, e_6, e_7\}$

Let present the edge set  $E_1$  &  $E_2$  in graphical form

And



As we can see that the edge set  $E_2$  is containing all the edges of  $E_1$  is the subset of the edge set  $E_2$ . It can be shown as:-



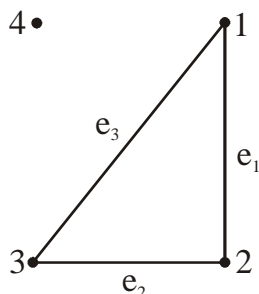
On the same time, we can say  $E_2$  is the superset of  $E_1$  as it contains all the edges of  $E_1$  i.e.  $E_2 \supset E_1$ .

**Proper Subset:** The edge set  $E_1$  is said to be proper subset of  $E_2$  if every edge of  $E_1$  is contained in the set  $E_2$  but there exist at least on edge of  $E_2$  which does not lie in the set  $E_1$ .

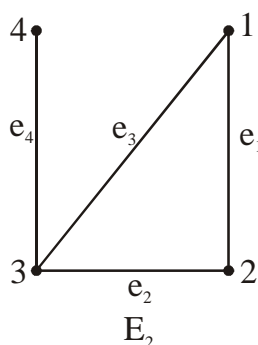
i.e.  $E_1 \subset E_2$  but  $E_1 \neq E_2$

**Example :** let us consider two edge set  $E_1$  and  $E_2$  such that

$E_1 = \{e_1, e_2, e_3\}$  i.e.



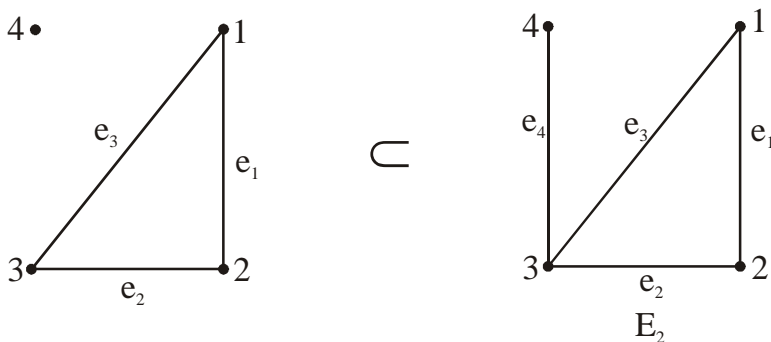
and the another edge set  $E_2$



can be represented as following:-

On analyzing, both the edge set, we can conclude that the edge set  $E_2$  contains all the edges of the edge set  $E_1$  and one more edge is there in  $E_2$  which is not in the edge set  $E_1$ . i.e.  $E_1 \neq E_2$

Thus  $E_1$  is the proper subset of  $E_2$ .



**Equivalent Sets :**

Two edge set  $E_1$  and  $E_2$  are said to be equivalent if the number of edges in  $E_1$  and  $E_2$  are same whether the edges are same or not i.e. the no. of edges should be same.

OR



two edge set are said to be equivalent if the number of edges are same irrespective of the fact that whether edges are same or not.

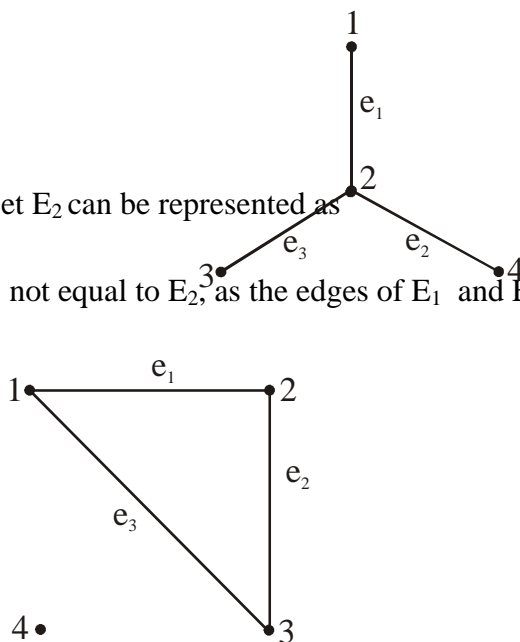
i.e.  $n(E_1) = n(E_2)$

**Example:** let us assume two edge set  $E_1$  and  $E_2$  such that:

$E_1 = \{e_1, e_2, e_3\}$  i.e.

and the another edge set  $E_2$  can be represented as

Now the edge set  $E_1$  is not equal to  $E_2$ , as the edges of  $E_1$  and  $E_2$  are not same.



But no. of edges are same, it mean there exist there edges in both the edge set.

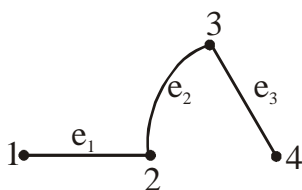
Hence  $E_1 \sim E_2$

**Power Set :**

If an edge set  $e$  is defined on  $n$  number of edges, then the power set of  $E$  will contain  $2^n$  elements i.e. all the subsets of  $E$  and denoted by  $P(E)$ .

let's verity it with the help of example.

**Example:** let us assume an edge set  $E$  on the given graph  $G$ .



As the given edge set  $E$  is containing three edges, the  $P(E)$  will contain  $2^3$  edge set i.e. 8 edge subset.

## 2. CONCLUSION

One more charming component of a chart based element determination is its possible versatility. As opposed to advance determination and in reverse end in old style relapse examination, the five analyzed component choice methodologies enjoy the benefit of being without model, as per the manner in which they are formed. Thus, diagram based strategies can be utilized as the pre-preparing period of computational or factual models within the sight of enormous scope or turbulent time series information in a more extensive scope of spaces.

Diagrams are huge in light of the fact that they are a visual portrayal of data. A chart portrays information that is what could be compared to many words. A chart can give data that is hard to depict in words. IN A WRITTEN LETTER TO C. G.W. Huygens, Huygens of 1679 Leibniz was disappointed with the typical organize math treatment of mathematical shapes, asserting that "one more sort of investigation, mathematical or straight, which manages position, as polynomial math does with greatness" . Actually, Leibniz was quick to explore the ageometry of positionso (geometria situs), which, as per L. "It is only worried about the assurance of area, and its ascribes; it doesn't need estimations or calculations finished with them," Euler said plainly in his popular 1736 Konigsberg spans article, which needed to address the introduction of chart hypothesis. With regards to software engineering, the utilization of chart hypothesis is outrageous. Chart hypothesis can promptly resolve numerous issues that are hard to decide or carry out. In diagram hypothesis, there are a wide range of kinds of charts. Each diagram type has its own arrangement of qualities. Most applications utilize one of these charts to adjust their answers for difficulties. Numerous issues can be addressed as charts and addressed effectively because of the portrayal strength and adaptability of diagrams. Asset designation, distance minimization, network building, ideal way distinguishing proof, information mining, circuit minimization, picture catching, and picture preparing are a portion of the issues that chart hypothesis can deal with.

Diagram hypothesis has turned into a somewhat critical subject in science because of the steady review done in the field. Various kinds of diagrams are remembered for chart hypothesis, every one of which has fundamental chart includes just as certain extra characteristics. These attributes recognize a diagram from different charts of a similar sort.

These components decide how a chart's vertex and edges are orchestrated in a given design. There are an assortment of activities that might be performed on different diagram types. Accordingly, diagram hypothesis is a major and requesting subject. Diagrams, then again, are utilized in a wide scope of uses as a solid instrument for addressing tremendous and complex issues. Diagrams can be utilized to answer issues in an assortment of spaces, including science, science, software engineering, and functional exploration. Subsequently, chart hypothesis is valuable in an assortment of uses, large numbers of which are broadly utilized in reality. Diagram hypothesis is utilized in pretty much every discipline these days, including search PC organizations. To effectively execute and deal with these applications, a strong comprehension of diagram hypothesis is required. Commonly, materials can't address all parts of chart hypothesis. Materials that accurately cover every single component of chart hypothesis miss the mark regarding giving a short outline of how those ideas are applied in certifiable circumstances. Chart hypothesis materials every now and again neglect to portray the basics of diagrams and their components. The creators of this paper look to give fundamental establishments of diagram hypothesis just as a careful comprehension of how these establishments are utilized.

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